



Title	計算世界観の深化と展開 : 東京工業大学グローバルCOEのご紹介
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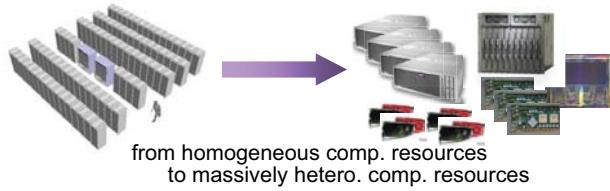
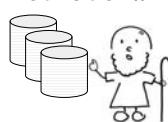
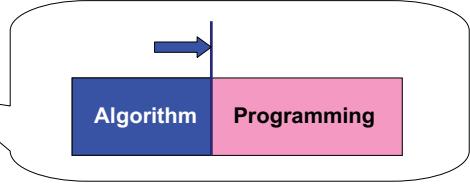
Proposal from Theory Group

ALGORITHMICS  
Algorithm Theory Research  
for Massively Hetero. Parallel Computing

algorithmics  $\leftrightarrow$  mathematics

topic (1) with Fukunaga, Kishimoto, Kobayashi

topic (2) with Okamoto, Onsjö

ALGORITHMICSfundamental  
Motivation#1. Paradigm shift in HPCMotivation#2. Heavy & hetero.  
computational demandsALGORITHMICSNew Algorithmic Research for  
Opt. Parallel Computing*Why new models ??*

- (1) Revisit of theoretical model: BSP, LogP
- (2) New model for new computing, e.g., GPU

ALGORITHMICS

東京工業大学 グローバルCOEのご紹介

## 計算世界観の深化と展開

## 拠点目標

教育：計算世界観の手法を実践できる人材の育成  
研究：計算を中心とした科学の新たな科学の確立

## 計算世界観とは

計算を中心に科学の対象を見よう！  
それにより新しい科学の手法を導こう！

計算 = 処理・判断の組み合わせ

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## Targets and Goals

## 1. How to handle mass. hetero parallelism

ultimate goal

Autonomic Computing

- run time self error detection  $\rightarrow$  self-healing
- self status recognition  $\rightarrow$  self-configuration

## 2. How to make use of mass. parallelism



ultimate goal

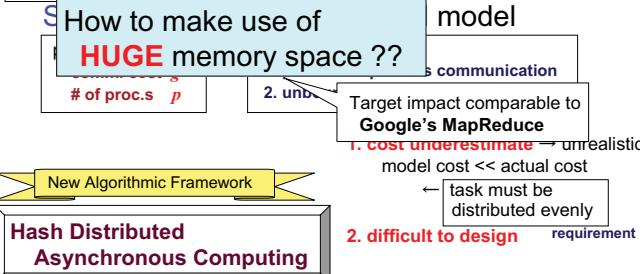
Opt. Parallel Computing

- (1) revisit of theoretical model: BSP, LogP
- (2) new model for new computing, e.g., GPU

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## Topic (1) Revisit of theoretical model

parallel comp. model

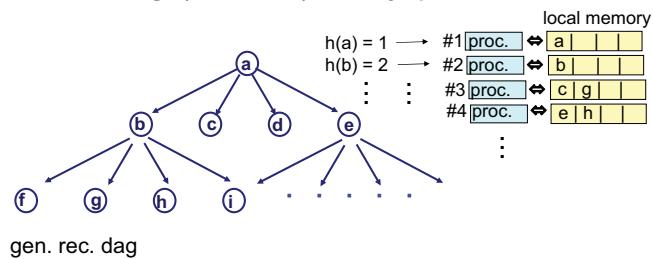


## ALGORITHMICs

### Extend Hash Distributed Computing

New Generic Parallel Algorithm for Dynamic Prog.

- asynchronous parallelism
- use huge (distributed) memory space



## ALGORITHMICs

### Topic (2) New model for new computing

GPU computing

my understanding  
(based on Tesla)

global memory

access cost  $T_{ext}$  ↕ block size  $B$

shared memory (cache)  
shared memory size  $M$

local processors  
warp size  
# of processors  $W$

access cost  $T_{sh}$   
to the shared memory  
no bank conflict case

Points of interest:

- + very fast local memory (cache)
- + local parallelism by local proc.s
- + stream processing (?)
- bank conflict ⇒ slow down