Craig Interpolation for a Sequent Calculus for Combining Intuitionistic and Classical Propositional Logic

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1 Introduction

This paper establishes the Craig interpolation for a multi-succedent sequent calculus for a combination of intuitionistic and classical propositional logic, denoted by $\mathsf{G}(\mathbf{C}+\mathbf{J})$. The calculus was provided in [16] and is based on the semantics offered in [4, 5]. The logic, called $\mathbf{C}+\mathbf{J}$, has two implications: intuitionistic and classical one[1]. They are interpreted in the Kripke semantics as follows (cf. [4, 5]):

$$\begin{array}{lll} w \models_M A \to_{\mathtt{i}} B & \text{iff} & \text{for all } v \in W, (wRv \text{ and } v \models_M A \text{ jointly imply } v \models_M B), \\ w \models_M A \to_{\mathtt{c}} B & \text{iff} & w \models_M A \text{ implies } w \models_M B, \end{array}$$

where M is an intuitionistic Kripke model, w is a possible world in M, and R is a preorder equipped in M. However this semantic treatment breaks one feature of intuitionistic logic called *heredity*, which is defined as: $w \models A$ and wRv jointly imply $v \models A$ for all Kripke models M and all states w and v in M. It is a well-known fact that this feature corresponds to an intuitionistically valid formula $A \to_{\mathtt{i}} (B \to_{\mathtt{i}} A)$. Therefore, the formula is not valid in the Kripke semantics of $\mathbf{C}+\mathbf{J}$. In order to avoid the formula being derivable in $\mathsf{G}(\mathbf{C}+\mathbf{J})$, the right rule for the intuitionistic implication should be restricted as follows:

$$\frac{A, C_1 \to_{\mathtt{i}} D_1, \ldots, C_m \to_{\mathtt{i}} D_m, p_1, \ldots, p_n \Rightarrow B}{C_1 \to_{\mathtt{i}} D_1, \ldots, C_m \to_{\mathtt{i}} D_m, p_1, \ldots, p_n \Rightarrow A \to_{\mathtt{i}} B} \ (\Rightarrow \to_{\mathtt{i}}).$$

The resulting calculus is sound and complete and a conservative extension of both an intuitionistic and a classical propositional sequent calculus (see [16]).

It is well-known that classical propositional logic and intuitionistic propositional logic enjoy the Craig interpolation theorem:

> If $A \to B$ is derivable, then there exists a formula C such that both $\Rightarrow A \to C$ and $\Rightarrow C \to B$ are also derivable and that $\mathsf{Prop}(C) \subseteq \mathsf{Prop}(A) \cap \mathsf{Prop}(B)$,

where $\mathsf{Prop}(D)$ denotes the set of all propositional variables in a formula D. The theorem can be shown in terms of a classical sequent calculus $\mathbf{LK}$ by Maehara's method in [9]. In multi-succedent intuitionistic sequent calculus $\mathbf{mLJ}$, the theorem can also be shown, though some modification of the ways is needed, as is noted in [10]. Since $\mathbf{C}+\mathbf{J}$ contains the two kinds of implication, the two types of Craig interpolation theorem can be considered in $\mathsf{G}(\mathbf{C}+\mathbf{J})$.

[1] In addition to $\mathbf{C}+\mathbf{J}$, other attempts to combine intuitionistic and classical logic are displayed in [1, 2, 3, 6, 7, 11, 12, 13, 14].

2 Syntax, Kripke Semantics and Sequent Calculus

2.1 Syntax and Kripke Semantics

This section reviews the syntax and the Kripke semantics of $\mathbf{C}+\mathbf{J}$. The syntax is defined in [16], and the Kripke semantics is based on the ones in [4, 5]. The syntax $\mathcal{L}$ consists of a countably infinite set Prop of propositional variables and the following logical connectives: falsum $\bot$, disjunction $\vee$, conjunction $\wedge$, intuitionistic implication $\to_{\mathtt{i}}$, and classical implication $\to_{\mathtt{c}}$. The set Form of all formulas in our syntax is defined inductively as follows:

$$A ::= p \mid \bot \mid A \vee A \mid A \wedge A \mid A \to_{\mathtt{i}} A \mid A \to_{\mathtt{c}} A,$$

where $p \in \mathsf{Prop}$. We define $\top := \bot \to_{\mathtt{i}} \bot$, $\neg_{\mathtt{c}} A := A \to_{\mathtt{c}} \bot$ and $\neg_{\mathtt{i}} A := A \to_{\mathtt{i}} \bot$.

Let us move to the semantics for the syntax $\mathcal{L}$.

Definition 1. A *model* is a tuple $M = (W, R, V)$ where

- W is a non-empty set of possible worlds,
- R is a preorder on W, i.e., R satisfies reflexivity and transitivity,
- $V : \mathsf{Prop} \to \mathcal{P}(W)$ is a valuation function satisfying the following *heredity* condition: $w \in V(p)$ and wRv jointly imply $v \in V(p)$ for all worlds $w, v \in W$.

Definition 2. Given a model $M = (W, R, V)$, a world $w \in W$ and a formula A, the *satisfaction relation* $w \models_M A$ is inductively defined as follows:

$$\begin{array}{lll} w \models_M p & \text{iff} & w \in V(p), \\ w \not\models_M \bot, & & \\ w \models_M A \wedge B & \text{iff} & w \models_M A \text{ and } w \models_M B, \\ w \models_M A \vee B & \text{iff} & w \models_M A \text{ or } w \models_M B, \\ w \models_M A \to_{\mathtt{i}} B & \text{iff} & \text{for all } v \in W, (wRv \text{ and } v \models_M A \text{ jointly imply } v \models_M B). \\ w \models_M A \to_{\mathtt{c}} B & \text{iff} & w \models_M A \text{ implies } w \models_M B. \end{array}$$

Let us say that a formula A is a *semantic consequence* of a set of formulas Γ, represented as $\Gamma \models A$, if $w \models_M C$ for any formula $C \in \Gamma$, then $w \models_M A$ for all models $M = (W, R, V)$ and all worlds $w \in W$. We use $\Gamma \models \Delta$ if $\Gamma \models A$ for some formula $A \in \Delta$. We say that A is *valid* if $\emptyset \models A$ holds. We say a formula A satisfies *heredity* if the following holds: $w \models A$ and wRv jointly imply $v \models A$ for all Kripke models M and all states w and v in M.

Proposition 1. A formula $\neg_{\mathtt{c}} p$ does not satisfy heredity.

Proposition 2. Neither $\neg_{\mathtt{c}} p \to_{\mathtt{i}} (\top \to_{\mathtt{i}} \neg_{\mathtt{c}} p)$ nor $\neg_{\mathtt{c}} p \to_{\mathtt{c}} (\top \to_{\mathtt{i}} \neg_{\mathtt{c}} p)$ is valid.

Proposition 2 implies that an intuitionistic tautology $A \to_{\mathtt{i}} (B \to_{\mathtt{i}} A)$, which is known for the correspondence to heredity in intuitionistic logic, is no longer valid.

2.2 Multi-succedent sequent calculus $\mathsf{G}(\mathbf{C}+\mathbf{J})$

This section reviews the sequent calculus $\mathsf{G}(\mathbf{C}+\mathbf{J})$ provided in [16]. In what follows, we use the ordinary notion of multi-succedent sequent. A *sequent* is a pair of finite multisets denoted by $\Gamma \Rightarrow \Delta$, which is read as "if all formulas in Γ hold then some formulas in Δ hold." Table 1 provides our multi-succedent sequent calculus $\mathsf{G}(\mathbf{C}+\mathbf{J})$, where the notion of derivability is defined as an existence of a

Table 1: Sequent Calculus $\mathsf{G}(\mathbf{C}+\mathbf{J})$

Axioms

$$\frac{}{A \Rightarrow A}\,(Id) \quad \frac{}{\bot \Rightarrow}\,(\bot)$$

Structural Rules

$$\frac{\Gamma \Rightarrow \Delta}{\Gamma \Rightarrow \Delta, A}\,(\Rightarrow w) \quad \frac{\Gamma \Rightarrow \Delta}{A, \Gamma \Rightarrow \Delta}\,(w \Rightarrow) \quad \frac{\Gamma \Rightarrow \Delta, A, A}{\Gamma \Rightarrow \Delta, A}\,(\Rightarrow c) \quad \frac{A, A, \Gamma \Rightarrow \Delta}{A, \Gamma \Rightarrow \Delta}\,(c \Rightarrow)$$

$$\frac{\Gamma \Rightarrow \Delta, A \quad A, \Pi \Rightarrow \Sigma}{\Gamma, \Pi \Rightarrow \Delta, \Sigma}\,(Cut)$$

Propositional Logical Rules

$$\frac{A, C_1 \to_{\mathtt{i}} D_1, \ldots, C_m \to_{\mathtt{i}} D_m, p_1, \ldots, p_n \Rightarrow B}{C_1 \to_{\mathtt{i}} D_1, \ldots, C_m \to_{\mathtt{i}} D_m, p_1, \ldots, p_n \Rightarrow A \to_{\mathtt{i}} B}\,(\Rightarrow \to_{\mathtt{i}}) \quad \frac{\Gamma_1 \Rightarrow \Delta_1, A \quad B, \Gamma_2 \Rightarrow \Delta_2}{A \to_i B, \Gamma_1, \Gamma_2 \Rightarrow \Delta_1, \Delta_2}\,(\to_i \Rightarrow)$$

$$\frac{A, \Gamma \Rightarrow \Delta, B}{\Gamma \Rightarrow \Delta, A \to_{\mathtt{c}} B}\,(\Rightarrow \to_{\mathtt{c}}) \quad \frac{\Gamma_1 \Rightarrow \Delta_1, A \quad B, \Gamma_2 \Rightarrow \Delta_2}{A \to_{\mathtt{c}} B, \Gamma_1, \Gamma_2 \Rightarrow \Delta_1, \Delta_2}\,(\to_{\mathtt{c}} \Rightarrow)$$

$$\frac{\Gamma \Rightarrow \Delta, A \quad \Gamma \Rightarrow \Delta, B}{\Gamma \Rightarrow \Delta, A \wedge B}\,(\Rightarrow \wedge) \quad \frac{A, \Gamma \Rightarrow \Delta}{A \wedge B, \Gamma \Rightarrow \Delta}\,(\wedge \Rightarrow_1) \quad \frac{B, \Gamma \Rightarrow \Delta}{A \wedge B, \Gamma \Rightarrow \Delta}\,(\wedge \Rightarrow_2)$$

$$\frac{\Gamma \Rightarrow \Delta, A}{\Gamma \Rightarrow \Delta, A \vee B}\,(\Rightarrow \vee_1) \quad \frac{\Gamma \Rightarrow \Delta, B}{\Gamma \Rightarrow \Delta, A \vee B}\,(\Rightarrow \vee_2) \quad \frac{A, \Gamma \Rightarrow \Delta \quad B, \Gamma \Rightarrow \Delta}{A \vee B, \Gamma \Rightarrow \Delta}\,(\vee \Rightarrow)$$

finite tree, which is called a *derivation*, generated by inference rules of Table 1 from initial sequents (Id) and $(\bot)$ of Table 1.

Our basic strategy of constructing $\mathsf{G}(\mathbf{C}+\mathbf{J})$ is to add classical implication to the propositional fragment of multi-succedent sequent calculus $\mathbf{mLJ}$ of intuitionistic propositional logic, proposed by Maehara [8]. However, if the ordinary left and right rules of classical implication were added, the soundness of the resulting calculus would fail, because a formula $\neg_{\mathtt{c}} p \to_{\mathtt{c}} (\top \to_{\mathtt{i}} \neg_{\mathtt{c}} p)$, which is not valid by Proposition 2, would be derivable. This is the reason why the original right rule

$$\frac{A, \Gamma \Rightarrow B}{\Gamma \Rightarrow A \to_{\mathtt{i}} B}$$

of intuitionistic implication of $\mathbf{mLJ}$ is restricted to the right rule given in Table 1. Based on the abbreviation defined in Section 2.1, the following rules for negations are obtained respectively:

$$\frac{A, C_1 \to_{\mathtt{i}} D_1, \ldots, C_m \to_{\mathtt{i}} D_m, p_1, \ldots, p_n \Rightarrow}{C_1 \to_{\mathtt{i}} D_1, \ldots, C_m \to_{\mathtt{i}} D_m, p_1, \ldots, p_n \Rightarrow \neg_{\mathtt{i}} A}\,(\Rightarrow \neg_{\mathtt{i}}) \quad \frac{\Gamma \Rightarrow \Delta, A}{\neg_{\mathtt{i}} A, \Gamma \Rightarrow \Delta}\,(\neg_i \Rightarrow)$$

$$\frac{\Gamma, A \Rightarrow \Delta}{\Gamma \Rightarrow \neg_{\mathtt{c}} A, \Delta}\,(\Rightarrow \neg_{\mathtt{c}}) \quad \frac{\Gamma \Rightarrow \Delta, A}{\neg_{\mathtt{c}} A, \Gamma \Rightarrow \Delta}\,(\neg_{\mathtt{c}} \Rightarrow).$$

Proposition 3. For any $\Gamma \cup \Delta \subseteq \mathsf{Form}$, $\Gamma \Rightarrow \Delta$ is derivable in $\mathsf{G}(\mathbf{C}+\mathbf{J})$ iff $\Gamma \models \Delta$ holds.

Proposition 4. If $\Gamma \Rightarrow \Delta$ is derivable in $\mathsf{G}(\mathbf{C}+\mathbf{J})$, then $\Gamma \Rightarrow \Delta$ is derivable in $\mathsf{G}^{-}(\mathbf{C}+\mathbf{J})$, where $\mathsf{G}^{-}(\mathbf{C}+\mathbf{J})$ is the calculus obtained by removing the rule (Cut) from $\mathsf{G}(\mathbf{C}+\mathbf{J})$.

By Proposition 4, the subformula property is obtained, which ensures the calculus is a conservative extension of both intuitionistic and classical propositional logic.

3 Craig Interpolation

In this section, we establishes two types of Craig interpolation theorem for $\mathsf{G}(\mathbf{C}+\mathbf{J})$, based on Maehara's partition argument in [9]. This argument is originally for classical sequent calculus $\mathbf{LK}$, and is dependent on the fact that the cut elimination holds in the calculus. Since cut elimination holds also in $\mathsf{G}(\mathbf{C}+\mathbf{J})$, as is guaranteed by Proposition 4, this method can be employed. In the following part of this section, $\mathsf{Prop}(D)$ denotes the set of all propositional variables in a formula D. And if Γ is a finite multiset of formulas, we define $\mathsf{Prop}(\Gamma) = \bigcup\{\mathsf{Prop}(D) \mid D \in \Gamma\}$. Especially, we have $\mathsf{Prop}(\bot) = \emptyset$. We call $\langle(\Gamma_1 : \Delta_1); (\Gamma_2 : \Delta_2)\rangle$ a *partition* of a sequent $\Gamma \Rightarrow \Delta$, if Γ is Γ_1, Γ_2 and Δ is Δ_1, Δ_2. Let us say that C is an *interpolant* of $\langle(\Gamma_1 : \Delta_1); (\Gamma_2 : \Delta_2)\rangle$ if $\Gamma_1 \Rightarrow \Delta_1, C$ and $C, \Gamma_2 \Rightarrow \Delta_2$ are derivable and $\mathsf{Prop}(C) \subseteq \mathsf{Prop}(\Gamma_1, \Delta_1) \cap \mathsf{Prop}(\Gamma_2, \Delta_2)$.

Although the main idea of giving $\mathsf{G}(\mathbf{C}+\mathbf{J})$ is adding classical implication to intuitionistic logic, our proof is similar to that in classical logic. For establishing the Craig interpolation theorem for $\mathbf{mLJ}$, we cannot employ the notion of partition of the form $\langle(\Gamma_1 : \Delta_1); (\Gamma_2 : \Delta_2)\rangle$. This is because we cannot find an interpolant for $\langle(\emptyset : A); (A : \emptyset)\rangle$ as noted in [10]. Therefore, in order to show the theorem for $\mathbf{mLJ}$, the form of a partition should be restricted to $\langle(\Gamma_1 : \emptyset); (\Gamma_2 : \Delta)\rangle$. However, this restriction makes it possible to show neither of the two types of theorem in $\mathsf{G}(\mathbf{C}+\mathbf{J})$. Considering this situation, it seems difficult to establish the theorem for $\mathsf{G}(\mathbf{C}+\mathbf{J})$. However, the classical negation (or implication) enables us to use partitions of the form $\langle(\Gamma_1 : \Delta_1); (\Gamma_2 : \Delta_2)\rangle$ without any restriction to calculate an interpolant by Maehara method. This fact about the way of showing Craig interpolation theorem implies that $\mathbf{C}+\mathbf{J}$ can be regarded as the logic obtained by adding the special (intuitionistic) implication to classical logic[2].

Lemma 1. Suppose that $\Gamma \Rightarrow \Delta$ is derivable in $\mathsf{G}(\mathbf{C}+\mathbf{J})$. Then for any partition $\langle(\Gamma_1 : \Delta_1); (\Gamma_2 : \Delta_2)\rangle$ of the sequent, there exits an interpolant C in $\mathsf{G}(\mathbf{C}+\mathbf{J})$, i.e., such that both $\Gamma_1 \Rightarrow \Delta_1, C$ and $C, \Gamma_2 \Rightarrow \Delta_2$ are also derivable in $\mathsf{G}(\mathbf{C}+\mathbf{J})$, and $\mathsf{Prop}(C) \subseteq \mathsf{Prop}(\Gamma_1, \Delta_1) \cap \mathsf{Prop}(\Gamma_2, \Delta_2)$.

With Lemma 1, which is the core of the proof, we can easily show the following two types of Craig interpolation theorem.

Theorem 1. (Intuitionistic Craig Interpolation Theorem of $\mathsf{G}(\mathbf{C}+\mathbf{J})$). If $\Rightarrow A \to_{\mathtt{i}} B$ is derivable in $\mathsf{G}(\mathbf{C}+\mathbf{J})$, then there exists a formula C such that $\Rightarrow A \to_{\mathtt{i}} C$ and $\Rightarrow C \to_{\mathtt{i}} B$ are also derivable in $\mathsf{G}(\mathbf{C}+\mathbf{J})$ and that $\mathsf{Prop}(C) \subseteq \mathsf{Prop}(A) \cap \mathsf{Prop}(B)$.

Theorem 2. (Classical Craig Interpolation Theorem of $\mathsf{G}(\mathbf{C}+\mathbf{J})$). If $\Rightarrow A \to_{\mathtt{c}} B$ is derivable in $\mathsf{G}(\mathbf{C}+\mathbf{J})$, then there exists a formula C such that $\Rightarrow A \to_{\mathtt{c}} C$ and $\Rightarrow C \to_{\mathtt{c}} B$ are also derivable in $\mathsf{G}(\mathbf{C}+\mathbf{J})$ and that $\mathsf{Prop}(C) \subseteq \mathsf{Prop}(A) \cap \mathsf{Prop}(B)$.

4 Further Direction

In [15], the first-order expansion $\mathsf{G}(\mathbf{FOC}+\mathbf{J})$ of $\mathsf{G}(\mathbf{C}+\mathbf{J})$ can be given by adding classical universal quantifier to first-order multi-succedent intuitionistic sequent calculus $\mathbf{mLJ}$, although the similar restriction on the right rule for the intuitionistic universal quantifier is needed. Whether Craig interpolation holds in this expansion is an open question, which deserves being inquired.

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